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Presentation Content

- Trust as a business requirement for execution platforms
 - Result from the GST project
- Security Module Approach
 - Result from the GST project (contribution from KU Leuven)
- Partitioning Approach
 - Result from the MILS Project





Color Notation for Trust

Application Data	Communication Tunnel for Secure Data	
Application Data	Communication Tunnel for Confidential Data	
Application Data	Communication Tunnel for Authenticated Data	
Application Data	Communication Tunnel for Insecure Data	



Trust in an Execution Environment



 A node is structured into a « trusted part » and a « non trusted part »



- Very classical in an OS with privileged mode management
 - But Security is transversal







Case 1 of E2E Security

- A and B communicate via C
- C does not have access to application data
 - If C is malicious it can deny transmission







Case 2 or E2E Security

- A communicates with C, C communicates with B
- C has access to application data
 - it is a trusted node
- Crypto mechanisms can be different in A-C and in C-B







Case 3 of E2E Security

- A communicates with C, C communicates with B
 - Only a trusted part of C has access to application data, called security module
- Crypto mechanisms can be different in A-C and in C-B segments



Endpoints Might also be Secure





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Secure Execution Environment

Endpoints Must also be Secure





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Including Input-Outputs







Peer-to-peer Security == Circle of Trust?







Secure Execution Environment for V2V and V2I Communication



Vehicle TCU trusted part?

- TCU includes a trusted part
 - e.g. non trusted part is PC centric part of TCU, trusted part is CALM implementation part of TCU
- RSE trusted part?
- Control Center and Service Center trusted part?









Example

Secure messaging:

- Key agreement phase:
 - Ping pong messages
 - sendPing, receivePing, preparePong, receivePong
- After key agreement:
 - byte[] dataToShip=prepareForSend(SecurityLevel, Data, SessionAlias)
 - byte[] receivedData=processIncoming(incomingData, SecurityContext)
- Receiver engine:
 - Endless loop:
 - Message incomingData=receiveData()
 - Case(incomingData.type){
 - Ping: { Message pong=preparePong(ping);send(incomingData); }
 - Pong: { processPong(incomingData); }
 - Insecure: {...}
 - Confidential: {...}
 - Authenticated: {...}
 - Secure: {...}





Example

Secure data storage:

- storeData(SecurityLevel, Data, Alias, OverwriteIfExists)
 - SecurityLevel: plaintext, encrypted, integrity protected, confidential
 - Alias: (unique) reference to retrieve the data later on
 - OverwriteIfExists: self-explanatory boolean
- byte[] fetchedData=retrieveData(Alias, SecurityContext) throws noSuchAlias
 - SecurityContext: if the Alias refers to data which should not be made available given the current SecurityLevel, it will not successfully be fetched





Common Device Components







Examples of Security Modules

- Hardware security module (most expensive)
 - Used for high-bandwidth communications, secure payments, etc.
- Smartcard, SecurID token, SIM card
 - Commonly used to provide strong user, service and device authentication
- Trusted platform module (TPM)
 - By default built into many new laptops and desktops
 - Lacks features necessary for GST, e.g., authentication of users, application data, etc.
 - TPM only authenticates the device
- Software key store (cheapest)
 - Cryptography-related data is stored in persistent memory (flash, magnetic,...)
 - Non-secure microcontroller operates on this data





Security Modules Form Factors

- Dedicated coprocessor
 - Pluggable (e.g., reader for smartcard/memory card, SIM lock for SIM card, socket for chip
 - Fixed, e.g., soldered secure microprocessor (similar to smartcard, TPM)
- Using the main processor for functionality, coprocessor for important processes (e.g., payable services)
- Using the main processor only
 - Software-only security
 - Privileged mode (e.g. Arm with TrustZone)



Example of Use for V2V and V2I Communication

- Car A wishes to exchange data with Car B
- Car A steps
 - Use the Security Module of A to authenticate data
 - Send the authenticated data to B
- Car B steps:
 - Use the Security Module of B to validate the authenticity of received data
 - If authentication is OK, B processes data





MILS

- Multiple Independent Levels of Security
- Security Architecture for Middleware
- Based on military classification of security levels classifications
 - TS: top secret
 - S: secret
 - C: classified
 - U: unclassified



High Assurance MILS Architecture





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Secure Execution Environment





MILS

3 independent layers:

- Partitioning kernel
 - Offers process separation, in space and time
 - Small footprint => easier certification
- MILS middleware layer
- MILS application layer
 - Implement own security policies using provided protected mechanisms





Protection Mechanisms

Data isolation

- Information in the state of one partition must not be accessible to other partitions
- Information flow
 - Only authorized communication between partitions can occur
- Periods processing
 - Sanitization of shared resources between context switches
- Damage limitation
 - Failure in one partition is contained, so it does not affect other partitions





Independent Components for V2V and V2I Communication

- Which kind of independence
- Which kind of protection
 - In a typical microcontroller, a thread have access to the whole memory
 - can read sensitive data
 - can modify sensitive data





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